

Intelligent Anti-Money Laundering Solution Based upon Novel **Community Detection in Massive Transaction Networks on Spark**



Introduction

- AML system: rule-based, lack of pattern recognition function Divide complex structures into meaningful subgroups and calculate the suspicious degree for each group
- Temporal-directed Louvain algorithm
- Implement on Spark GraphX platform

Design and implementation

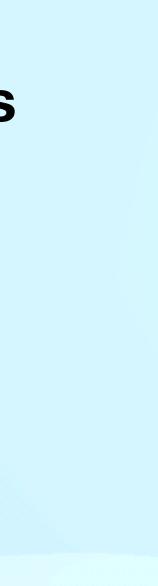
- A. Select effective maximal connected subgraph in massive transaction networks
- B. Community detection according to ML characteristics Edge weight optimization by node correction Temporal correction for edge weight Directed optimization for modularity
- C. Money laundering risk quantization for Communities

- Focus on the transaction data in a certain time period P_{t}
- Node: account, edge: transfer
- Merge the edge between node pairs with the same source and destination • $W_{Ri} = e^{\omega_m \cdot M_i + \omega_c \cdot C_i}$, where $\omega_m + \omega_c = 1$
- Divide the graph into different maximal connected subgraphs (MCS)
 - Filter the MCSs by the formula: V
- Hub nodes: the node whose degree exceeds a threshold D_{θ}
 - MCSs can be further filtered by Λ

A. Select effective maximal connected subgraph in massive transaction networks

$$V_{\theta 1} < V_{mcs} < V_{\theta 2}$$

$$N_{hubs} > N_{\theta}$$



A) Edge weight optimization by node correction

- Louvain algorithm: overlook the weight of node -> need corrections
- The suspicious degree of a node is also important
- Node correction for src: $\sigma_s = e^{\omega_{Mv} \cdot M_s + \omega_{Cv} \cdot C_s + \omega_{Dv} \cdot D_s}$
- Node correction for dst: $\sigma_d = e^{\omega_{Mv} \cdot M_d + \omega_{Cv} \cdot C_d + \omega_{Dv} \cdot D_d}$
- New edge weight: $W_{Ni} = \sigma_s * \sigma_d * W_{Bi}$

B. Community detection according to ML characteristics B) Temporal correction for edge weight

- t_{Δ}^{in} and t_{Δ}^{out} : node's average time point of all inbound and outbound transfers
- $T_{s \rightarrow d}$: average transfer time point for edge src -> dst
- Deg_s^{in} and Deg_s^{out} : the number of all inbound and outbound transfers for src





- **B.** Community detection according to ML characteristics B) Temporal correction for edge weight Pattern 1: "Centralized out after multi-transfer in" (for edge src \rightarrow dst)
- $Deg_s^{in} > Deg_s^{out}$ and $T_{s \to d} > t_s^{in}$: promote edge weight
- $Deg_s^{in} > Deg_s^{out}$ but $T_{s \to d} < t_s^{in}$: reduce edge weight
- The correction factor $\theta_s = e^{\beta_s \cdot \tau_s^{out}}$, where $\beta_s = \frac{Deg_s^{in} - Deg_s^{out}}{Deg_s^{in} + Deg_s^{out}}, \tau_s^{out}$

• $\theta_s > 1$ if $Deg_s^{in} > Deg_s^{out}$ and $T_{s \to d}$

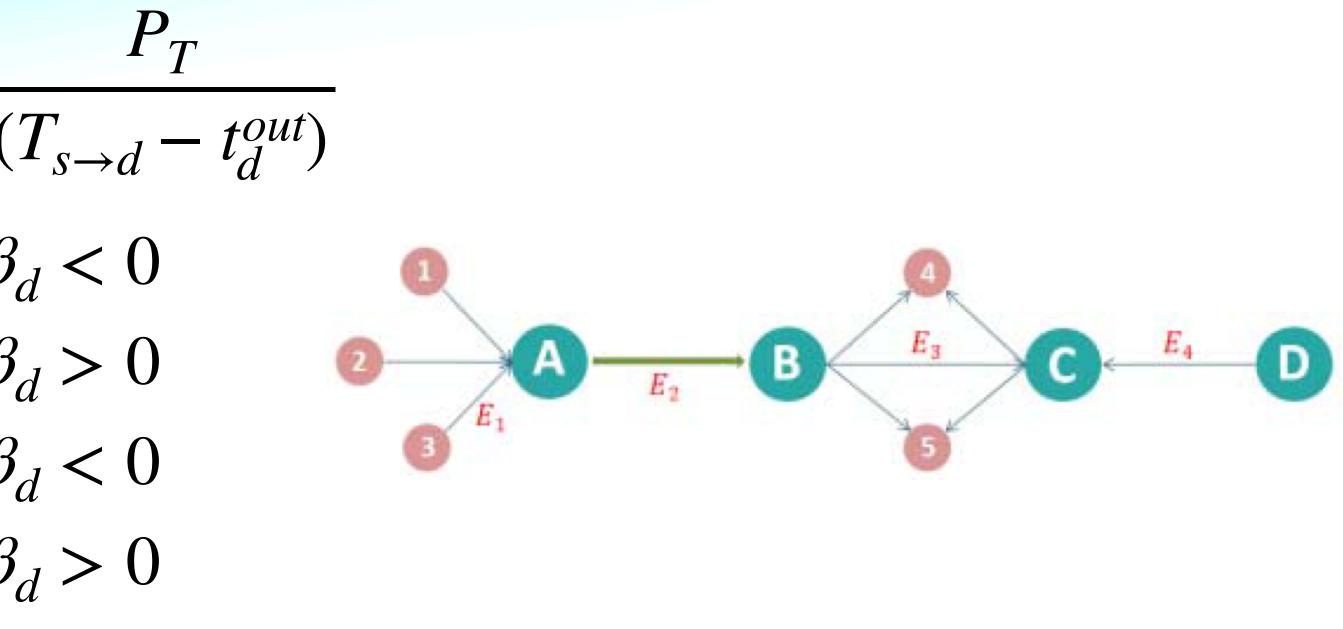
$$= \frac{P_T}{(T_{s \to d} - t_s^{in})}$$

$$t_s^{in}, \text{ else } \theta_s \leq t_s^{in}$$

B. Community detection according to ML characteristics B) Temporal correction for edge weight Pattern 2: "centralized in edge before multi-transfer out" (for destination node)

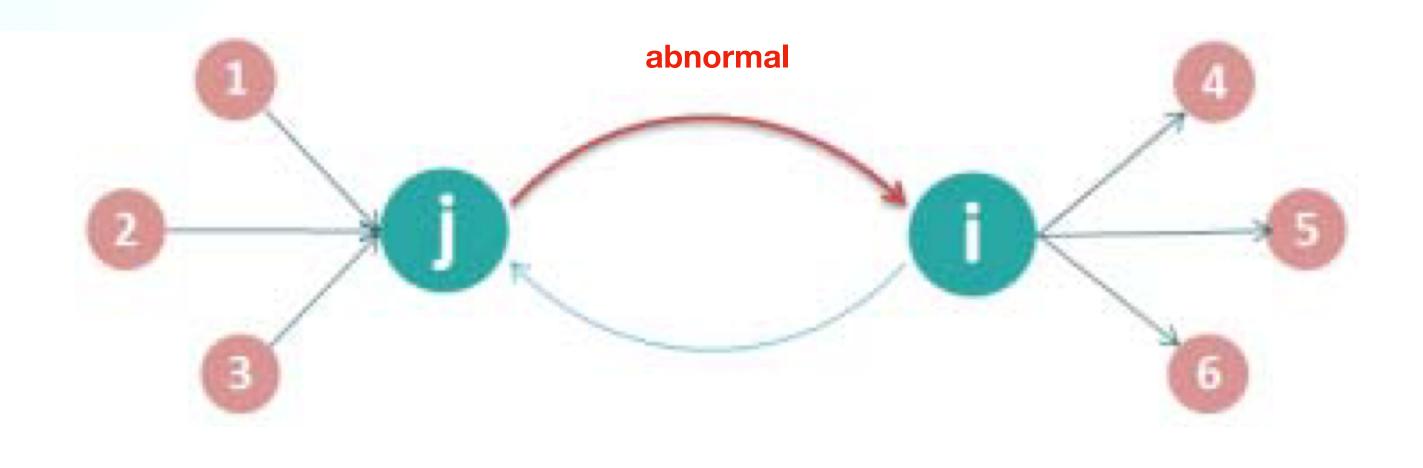
- $Deg_d^{int} < Deg_d^{out}$ and $T_{s \to d} < t_d^{out}$
- The correction factor $\theta_d = e^{\beta_d \cdot \tau_d^m}$, where $\beta_d = \frac{Deg_d^{in} - Deg_d^{out}}{Deg_d^{in} + Deg_d^{out}}$, $\tau_d^{in} = \frac{P_T}{(T_{s \to d} - t_d^{out})}$

$$W_{Ei} = \begin{cases} \theta_s \cdot \theta_d \cdot W_{Ni} , & \beta_s > 0 \text{ and } \beta_s \\ \theta_s \cdot W_{Ni} , & \beta_s > 0 \text{ and } \beta_s \\ \theta_d \cdot W_{Ni} , & \beta_s < 0 \text{ and } \beta_s \\ W_{Ni} , & \beta_s < 0 \text{ and } \beta_s \end{cases}$$



C) Directed optimization for modularity

- Example:
 - Node *i* : low in-degree, high out-degree
 - Node *j* : high in-degree, low out-degree



Consider the asymmetry of information caused by the direction of edges

C) Directed optimization for modularity

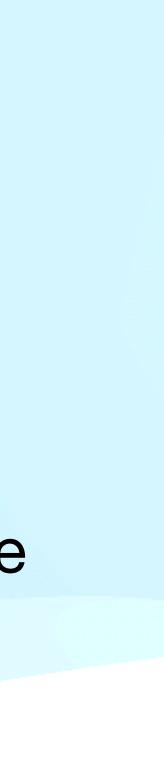
• Revision for Louvain algorithm:

• Revised factor $\delta_n = (k_n^{in} - k_n^{out})/k_n$

$$Q_D = \frac{1}{2m} \sum_{i,j} \left[A_{ij} - \frac{e^{\delta_i - \delta_j} k_i k_j}{2m} \right] \delta(c_i, c_j)$$
$$= \frac{1}{2m} \left[\sum_{i,j} A_{ij} - \frac{\sum_i e^{\delta_i} k_i \sum_j e^{-\delta_j} k_j}{2m} \right] \delta(c_i, c_j) = \sum_c \left[\frac{\sum W_{ec}}{2m} - \left(\frac{1}{2m} \right)^2 \sum M_c \right]$$

$$Q = rac{1}{2m} \sum_{ij} igg[A_{ij} - rac{k_i k_j}{2m} igg] \delta(c_i,c_j),$$

• Define k_n^{in} , k_n^{out} , k_n is the sum of edge weight linked in, out, and with node



- * Louvain algorithm:
 - Modularity: measures the relative density of edges inside communities with respect to edges outside communities L.L. 7 Г 1
 - Compare to a randomly rewired network
 - Divide an edge into 2 stubs -> m edges become 2m stubs
 - The probability for a stub of node *i* connecting with node *j*:

The probability for node i connecting with node j:

$$Q = \frac{1}{2m} \sum_{ij} \left[A_{ij} - \frac{\kappa_i \kappa_j}{2m} \right] \delta(c_i, c_j),$$

$$\frac{k_j}{2m-1}$$

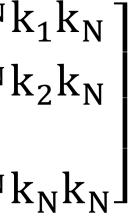
$$\frac{2m-1}{2m-1}$$

 k_{k}

B. Community detection according to ML characteristics • A_{ii} is the weight of the edge between i and j

- $m = \frac{1}{2} \sum_{i=1}^{n} A_{ij}$ is the sum of edge weights in the whole graph.
- $\delta(c_i, c_j) = 1$ if $c_i = c_j$, else $\delta(c_i, c_j) = 0$
- *c_i* is the community where node i belongs to
- k_i is the sum of all edge weights attached to node i
- M_c is the corresponding matrix for each community
- $\sum M_c$ is the sum of all elements in M_c
- means accumulating in the original community С

$$M_{c} = \begin{bmatrix} e^{\delta_{1} - \delta_{1}} k_{1} k_{1}, & e^{\delta_{1} - \delta_{2}} k_{1} k_{2}, & \cdots & e^{\delta_{1} - \delta_{N}} \\ e^{\delta_{2} - \delta_{1}} k_{2} k_{1}, & e^{\delta_{2} - \delta_{2}} k_{2} k_{2}, & \cdots & e^{\delta_{2} - \delta_{N}} \\ \vdots & \vdots & \ddots & \vdots \\ e^{\delta_{N} - \delta_{1}} k_{N} k_{1}, & e^{\delta_{N} - \delta_{2}} k_{N} k_{2}, & \cdots & e^{\delta_{N} - \delta_{N}} \end{bmatrix}$$



- Temporal-directed Louvain algorithm:
- Step 1: Initialize the community tag for each node by using its own node tag.

the allocation. Following is the revised formula for ΔQ_D

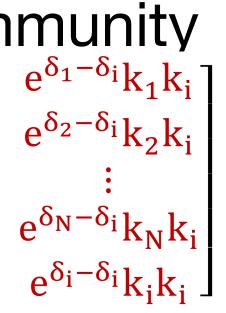
$$\Delta Q_{D} = \left[\frac{\sum W_{ec} + k_{i}^{c}}{2m} - \frac{1}{(2m)^{2}} \sum M_{cnew}\right] - \left[\frac{\sum W_{ec}}{2m} - \frac{1}{(2m)^{2}} \sum M_{c} - \frac{e^{\delta_{i} - \delta_{i}} k_{i} k_{i}}{(2m)^{2}}\right] = \frac{k_{i}^{C}}{2m} - \frac{\sum M_{cnew} - \sum M_{c} - k_{i} k_{i}}{(2m)^{2}} = \frac{k_{i}^{C}}{2m} - \frac{\Theta_{i}}{(2m)^{2}} \qquad (4)$$

$$\Theta_{i} = k_{i}e^{\delta_{i}}\sum_{c}k_{j}e^{-\delta_{j}} + k_{i}e^{-\delta_{i}}\sum_{c}k_{j}e^{\delta_{j}}$$

Step 2: For each node, maximize the difference ΔQ_D between allocating it in the original community and the neighbor community \rightarrow if the maximum $\Delta Q_D > 0$, actualize

• k_i^c : the total weight of edges formed between node *i* and all nodes in the community $[e^{\delta_1 - \delta_1}k_1k_1, e^{\delta_1 - \delta_2}k_1k_2, \cdots e^{\delta_1 - \delta_N}k_1k_N, e^{\delta_1 - \delta_i}k_1k_i]$ $e^{\delta_2-\delta_1}k_2k_1$, $e^{\delta_2-\delta_2}k_2k_2$, \cdots $e^{\delta_2-\delta_N}k_2k_N$, ${}^{j} M_{\text{cnew}} = \begin{bmatrix} \vdots & \vdots & \ddots & \vdots \\ e^{\delta_{N} - \delta_{1}} k_{N} k_{1}, & e^{\delta_{N} - \delta_{2}} k_{N} k_{2}, & \ddots & e^{\delta_{N} - \delta_{N}} k_{N} k_{N}, \end{bmatrix}$ $e^{\delta_i - \delta_1} k_i k_1$, $e^{\delta_i - \delta_2} k_i k_2$, ... $e^{\delta_i - \delta_N} k_i k_N$,





Temporal-directed Louvain algorithm:
 Step 3: Repeat Step 2 until the community tag of all nodes does not change.
 Step 4: Compress nodes with the same community label into a new node.
 Original total edge weight → self-link weight of new node
 Step 5: Repeat Step 2 until the modularity of the whole graph does not change

C. Money laundering risk quantization for Communities

- V_k : total nodes number
- E_k : total edge number
- M_k : total money amount
- D_k : average node suspicious degree

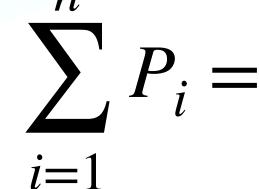
Temporal entropy $H_c = -\sum_{i}^{n} P_i * log_2 P_i$ (to measure the transfer time concentration)

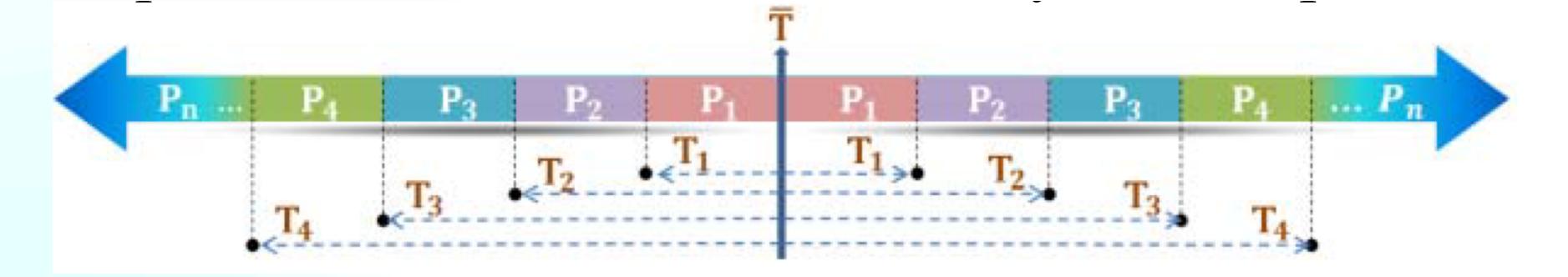
• ML risk score for the community k: $\Psi_k = e^{w_k \cdot V_k + w_E \cdot E_k + w_M \cdot M_k + w_D \cdot D_k + w_H \cdot H_k}$

C. Money laundering risk quantization for Communities

- \overline{T} : the average time point of a community
- ΔT : the absolute interval between each transfer time point and T
- P_1 is the percentage of transactions in the segment of $0 \le \Delta T \le T_1$
- P_2 is the percentage of transactions in the segment of $T_1 \leq \Delta T \leq T_2$

All segments should obey $\sum_{i=1}^{n} P_i = 1$



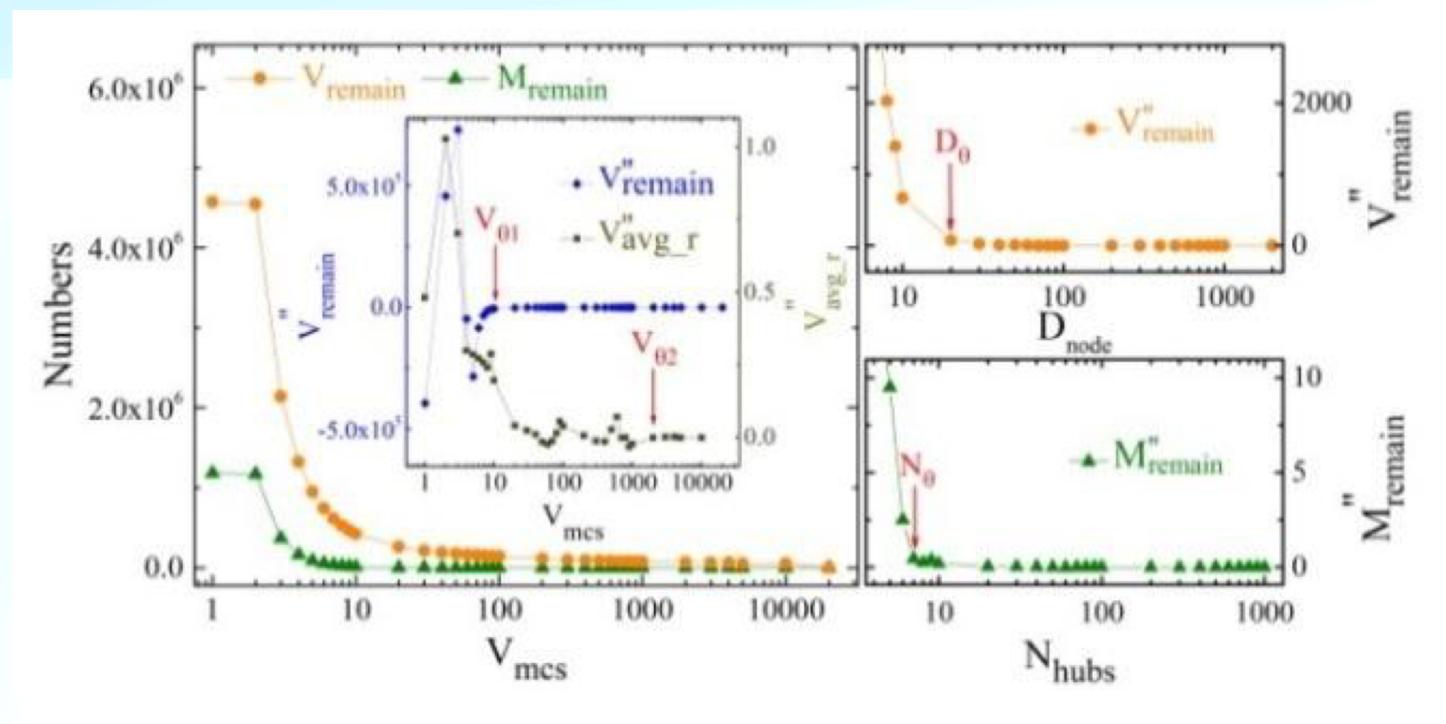


III. EXPERIMENTS AND RESULTS

A. Determination of Filter's Parameter

- 2016 -> filtered -> scale down to 45% of original size
- Filter's Parameter: $V_{\theta 1}$, $V_{\theta 2}$, D_{θ} , N_{θ}

•
$$V_{avg_r} = V_{remian} / M_{remain}$$



Data: around 10 million of real transfer records in the first week of November

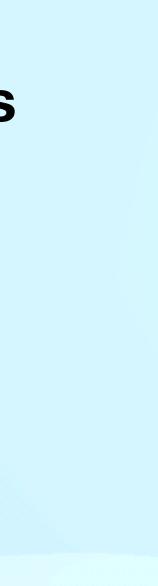


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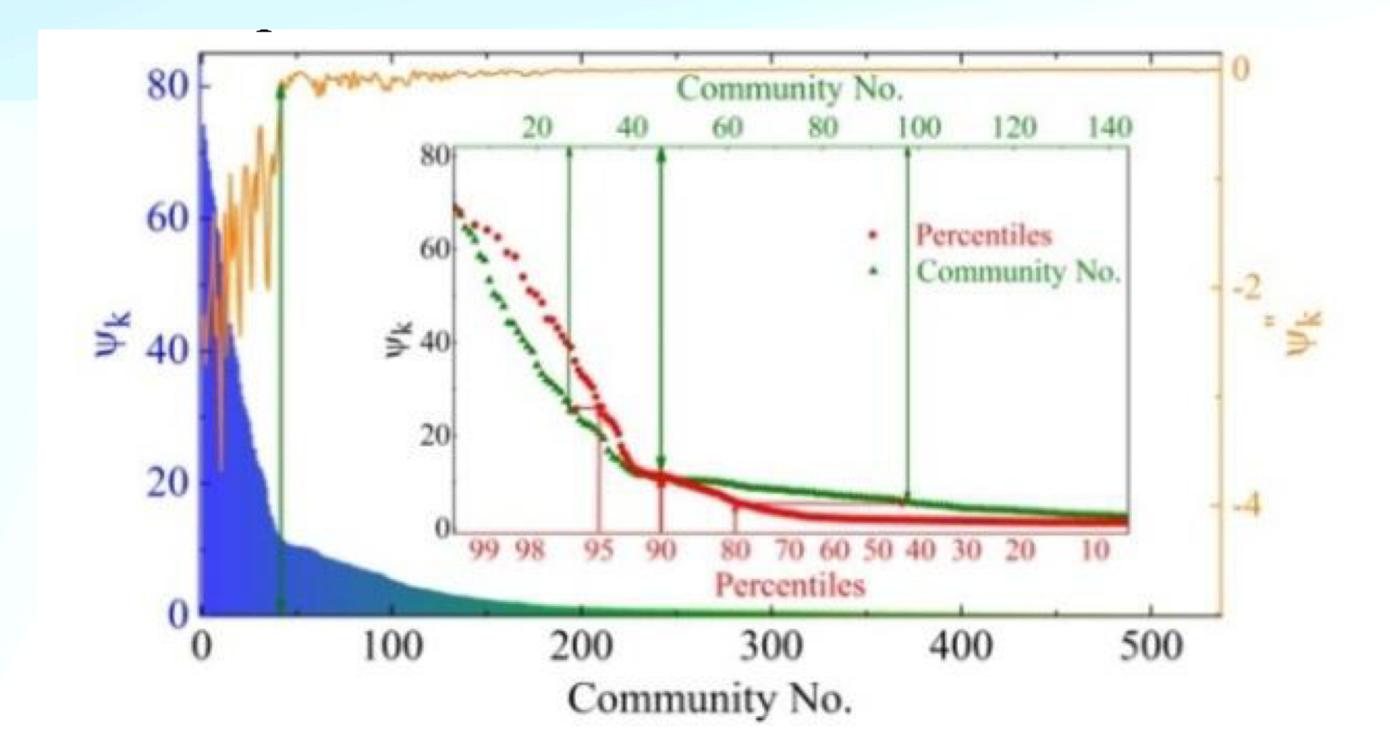
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B. ML risk level partition

- Orange line: first derivative for the rick score
- 95% ~ 100%: level 1, 90% ~ 95%: level 2, 80% ~ 90%: level 3
- After reporting level 1, 2 -> 9/13 were reconfirmed suspicious



C. TD Louvain algorithm v.s. Louvain algorithm

- Nodes ascribed to different communities are drawn in distinct colors
- Edges are also drawn in progressive color with increasing time

